# A Policy-free Kernel for Scalable Multi-core, Bounded Execution, and Verification

#### Gabe Parmer

Based on work by a village of strong researchers

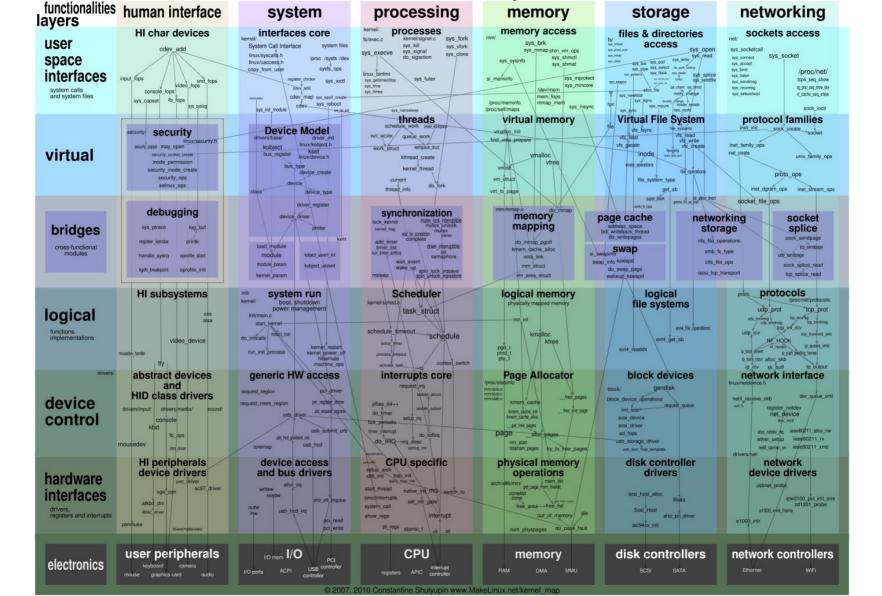


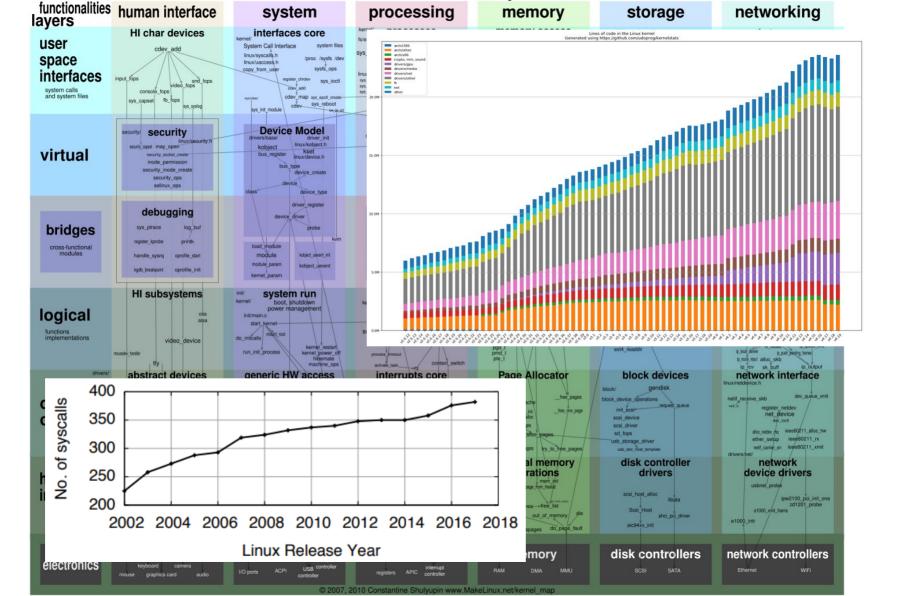
# System Design to Control Interference

#### Gabe Parmer

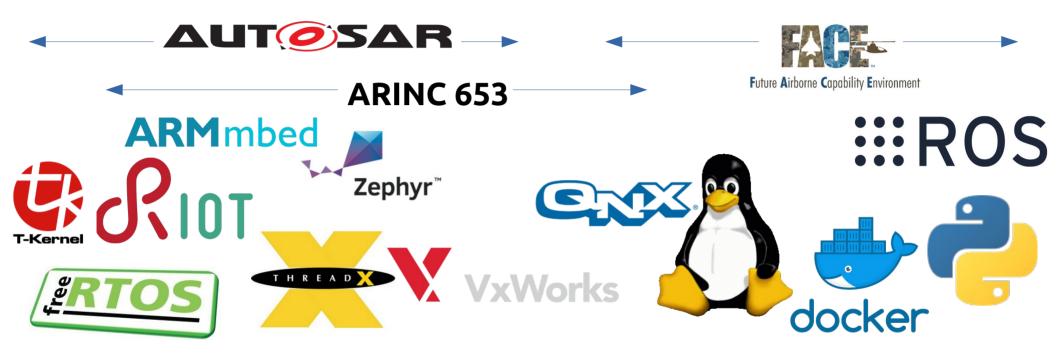
Based on work by a village of strong researchers







### Choose your trade-off



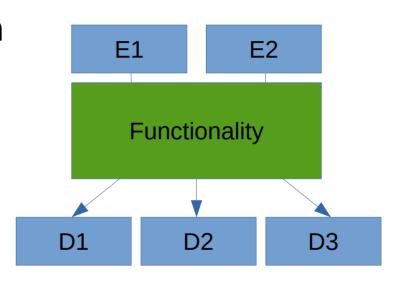
High Assurance Predictable **Simple**  Lower Assurance Tenuous predictability **Complex** 

### Goal:

#### Functionality + Dependability + Productivity

#### Component-based system design

- Code, data
- Export APIs (E1 = {fn, ...})
- Explicit dependencies (D1, ...)
- Unit of reuse & isolation

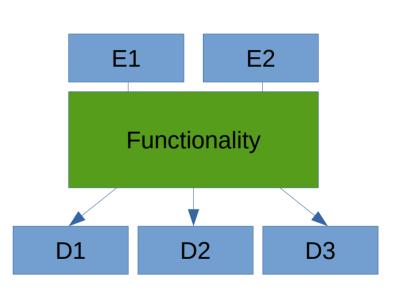


#### Goal:

#### Functionality + Dependability + Productivity

#### Component-based system design

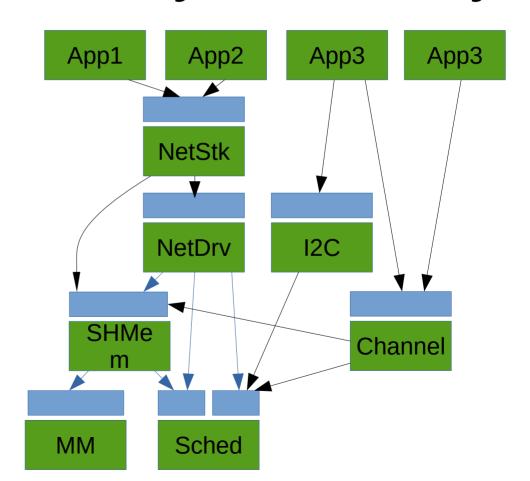
- · Code, data
- Export APIs (E1 = {fn, ...})
- Explicit dependencies (D1, ...)
- Unit of reuse & isolation
- Minimize functionality for the necessary APIs
- Strong, fine-grained isolation
- LCM & PoLP



# Goal: Functionality + Dependability + Productivity

 Components compose system

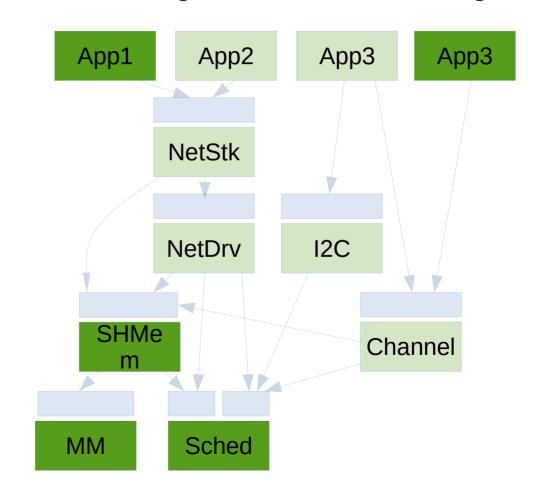
- Limit scope of
  - compromise
  - fault
  - unpredictability



# Goal: Functionality + Dependability + Productivity

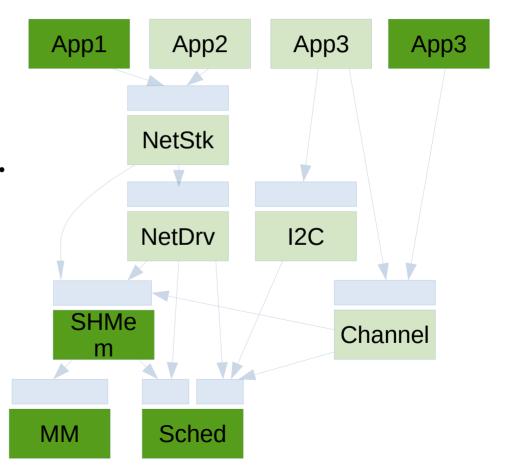
 Components compose system

- Limit *scope* of
  - compromise
  - fault
  - unpredictability



# Goal: Functionality + Dependability + Productivity

- Non-functional challenges
  - Shared resource interf.
  - Timing channels
  - \_

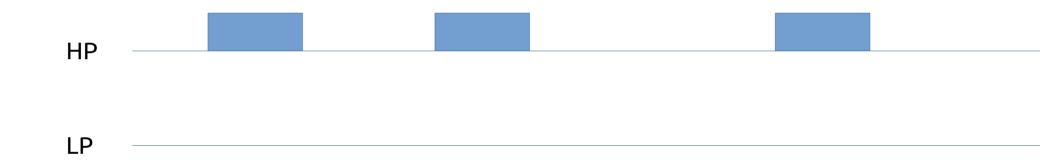


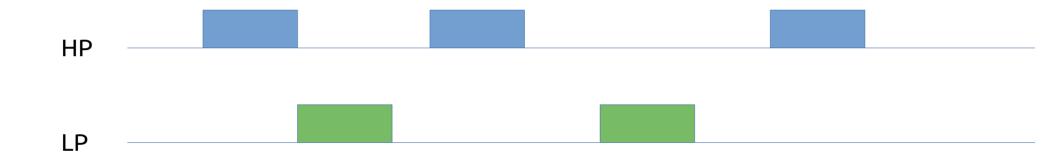
### Composite: OS for System Composition

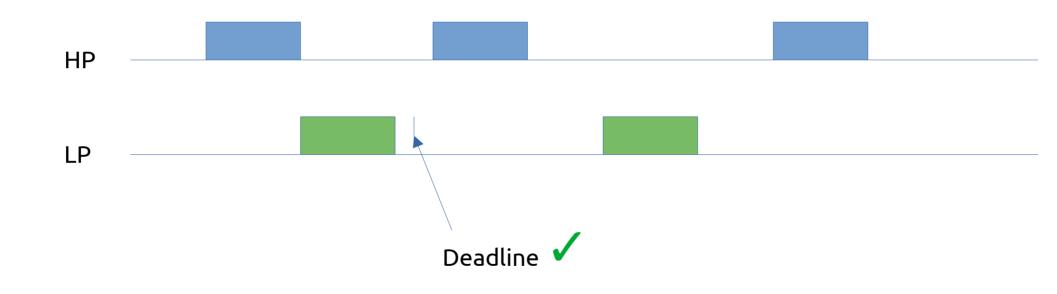
- Microkernel: small kernel, fast communication
  - User-level policies for resource mgmt
  - Page-table memory isolation between components
- Unique features get out of the way!
  - User-level scheduling → configurable policy
  - Wait-free kernel → predictably scalability
  - Integration with push-button verification

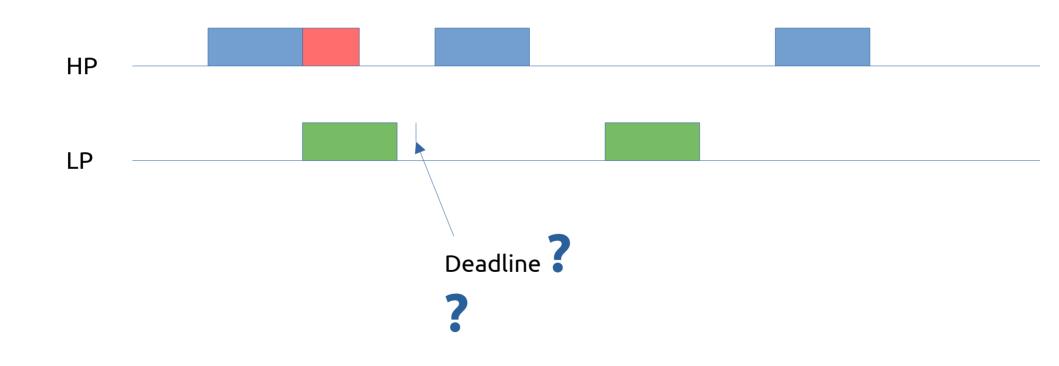
### OS Building Toolkit

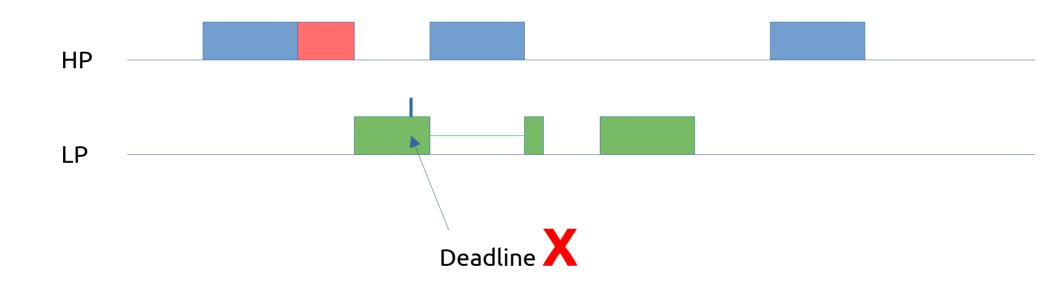
- Chaos managed interference for mixedcriticality systems
- EdgeOS per-client, multi-tenant isolation for edge processing
- Microsecond-scale multi-tenant infrastructure
- Multi-tenant host network functions
  - VM execution with NF interposition faster than baremetal!

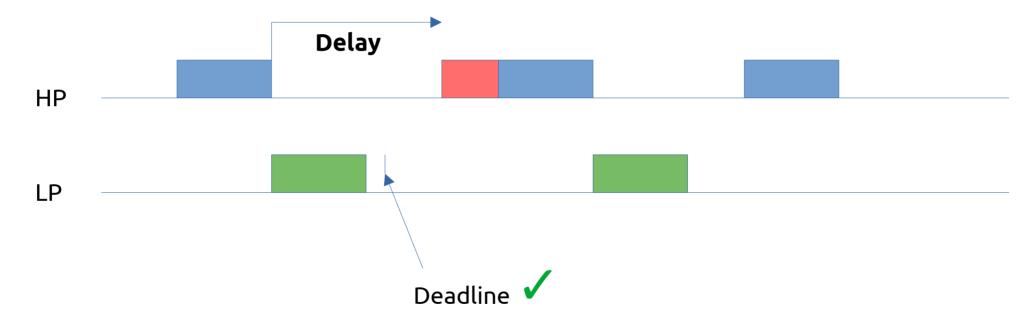


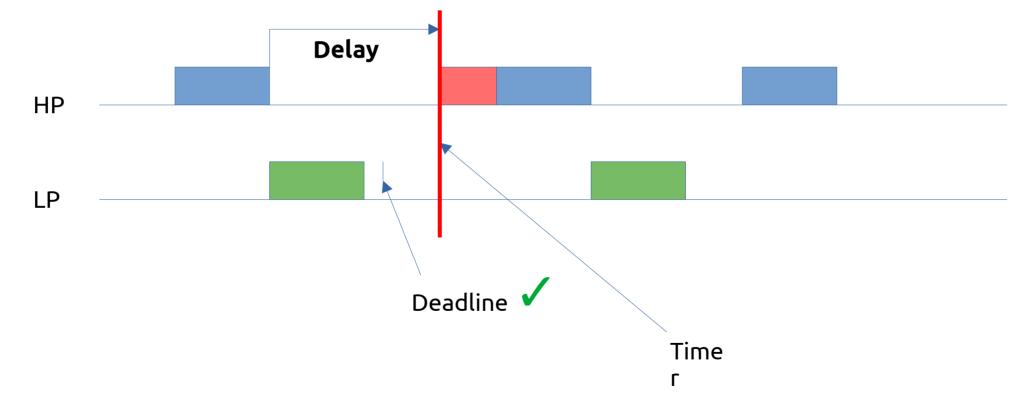












HP

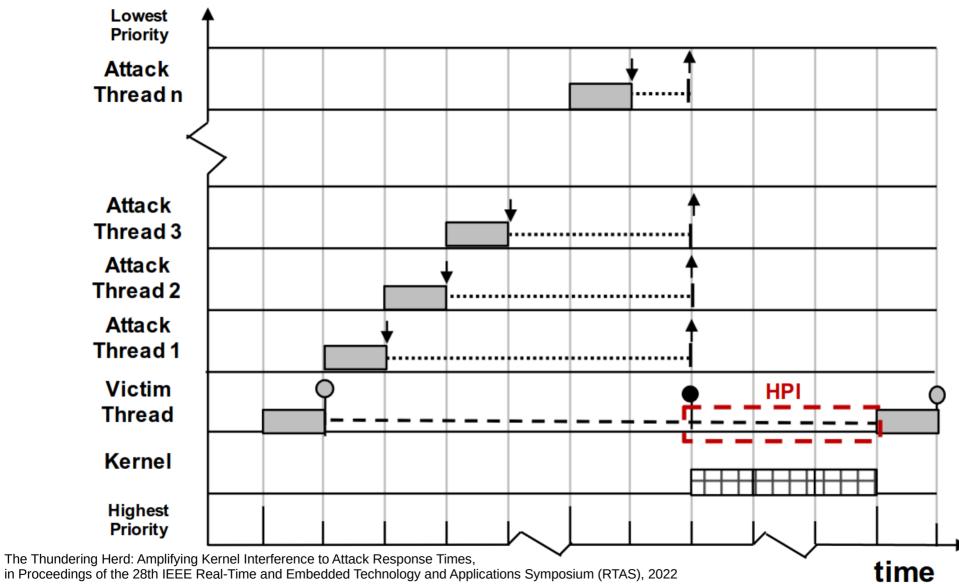
ΙP

Good policy:

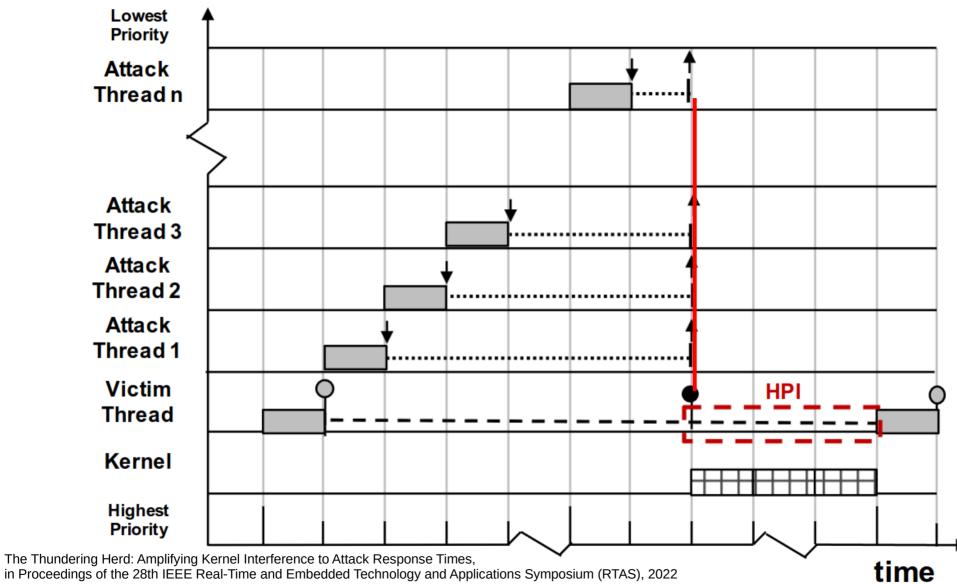
Limit impact of HP tasks on LP **Necessary for MCS** 

Risk:

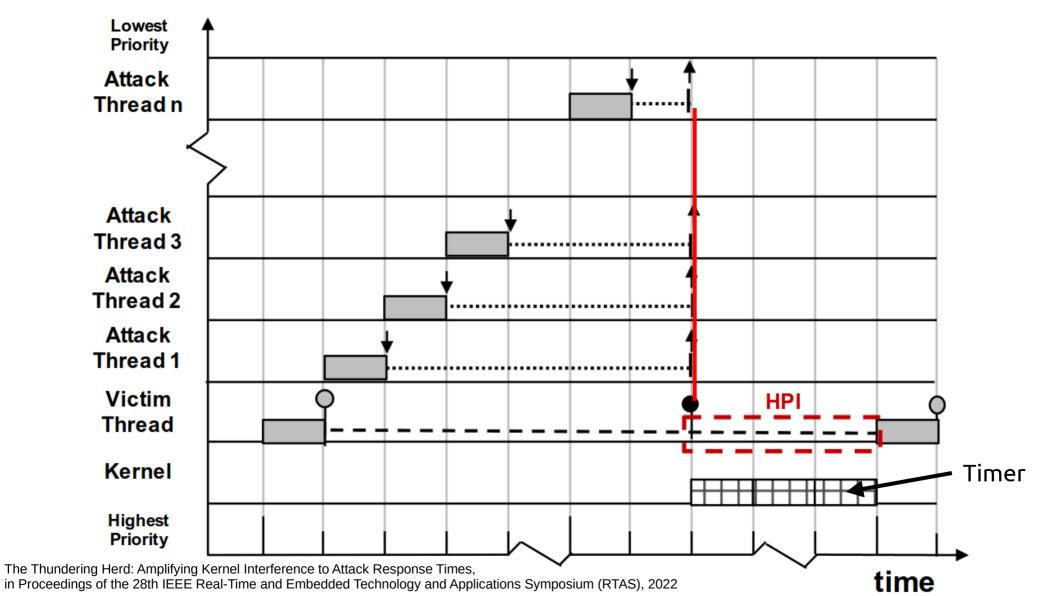
Policy becomes vector for attack

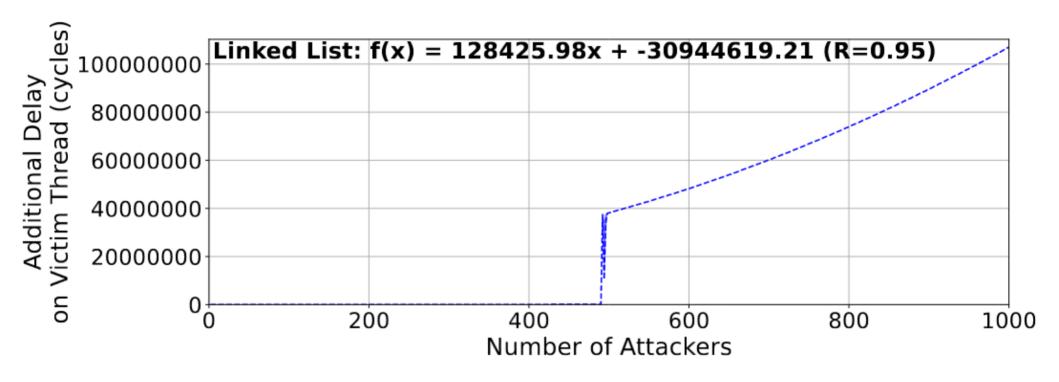


in Proceedings of the 28th IEEE Real-Time and Embedded Technology and Applications Symposium (RTAS), 2022



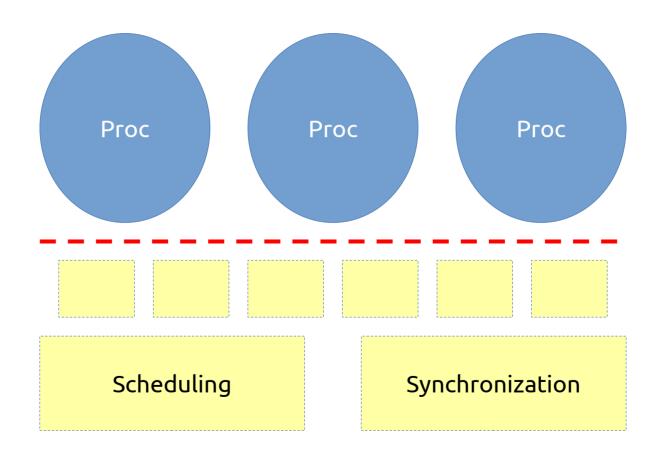
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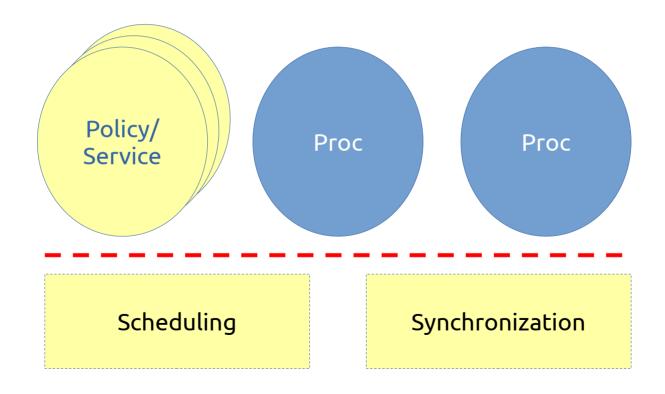




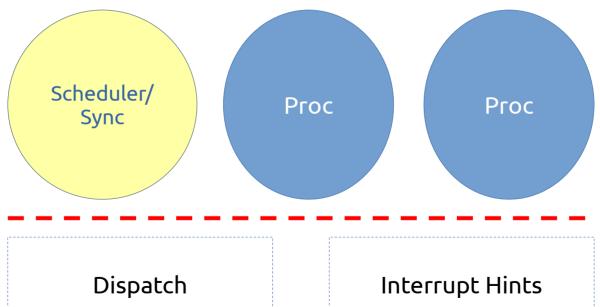
### Policy Freedom

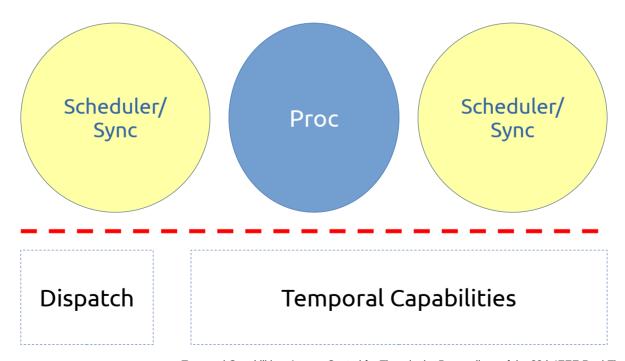
- Kernel policies are appealing
  - Simple
  - Centralized
- Policies will be attacked, when static
  - Policy impl. w/ interrupts disabled ...challenging





- Slite: OS Support for Near Zero-Cost, Configurable Scheduling, in Proceedings of the 26th IEEE Real-Time and Embedded Technology and Applications Symposium (RTAS), 2020
- HiRes: a System for Predictable Hierarchical Resource Management, to appear in Proceedings of the 17th IEEE Real-Time and Embedded Technology and Applications Symposium (RTAS 2011), Chicago, IL, April 2011
- "Predictable Interrupt Management and Scheduling in the Composite Component-based System", in Proceedings of the 29th IEEE Real-Time Systems Symposium (RTSS), Barcelona, Spain. 1-3 December 2008





Temporal Capabilities: Access Control for Time, in the Proceedings of the 38th IEEE Real-Time Systems Symposium (RTSS), 2017

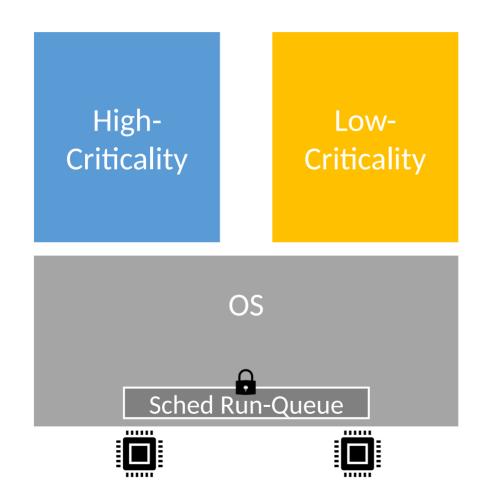
#### Kernel:

- No unbounded loops
- Actions: directly from syscall instructions
  - Capability-checked

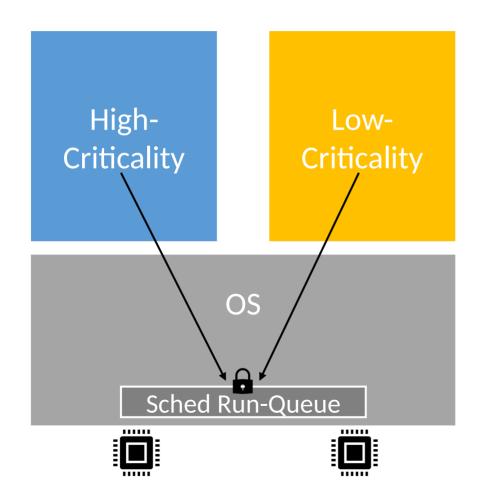
#### **Scheduler:**

- Any policy incl. emulating L4 IPC!
- Current work: constant-time everything

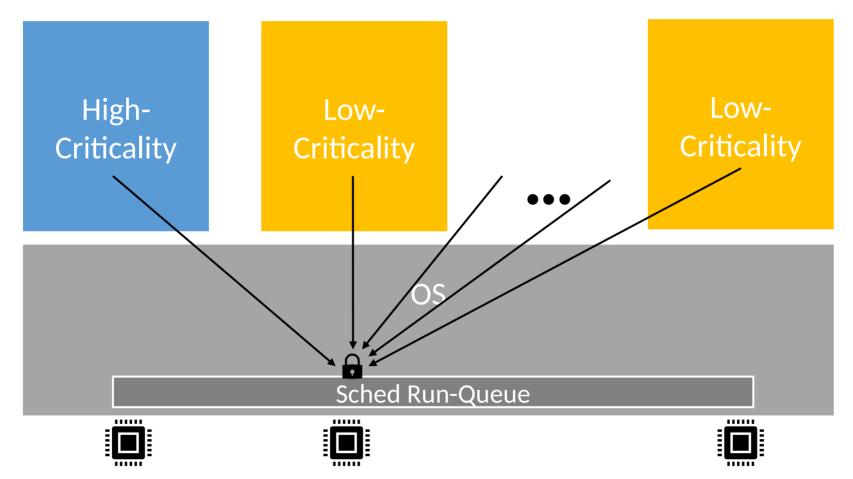
### Multi-Core: **Shared-Memory** Impact



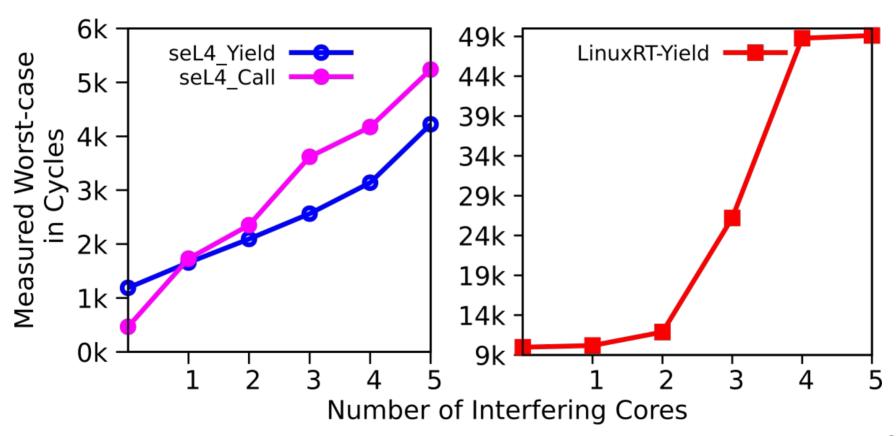
### Multi-Core: Shared-Memory Impact



### Multi-Core: **Shared-Memory** Impact



# Impact of Shared-Memory



### Composite µ-kernel

- Wait-free kernel data-structures
- No locks: all synchronization:
  - Atomic instructions
  - Without loops
  - Scalable memory reclamation using *time*
- Bounded execution!
  - → parallel kernel with *no locks!*

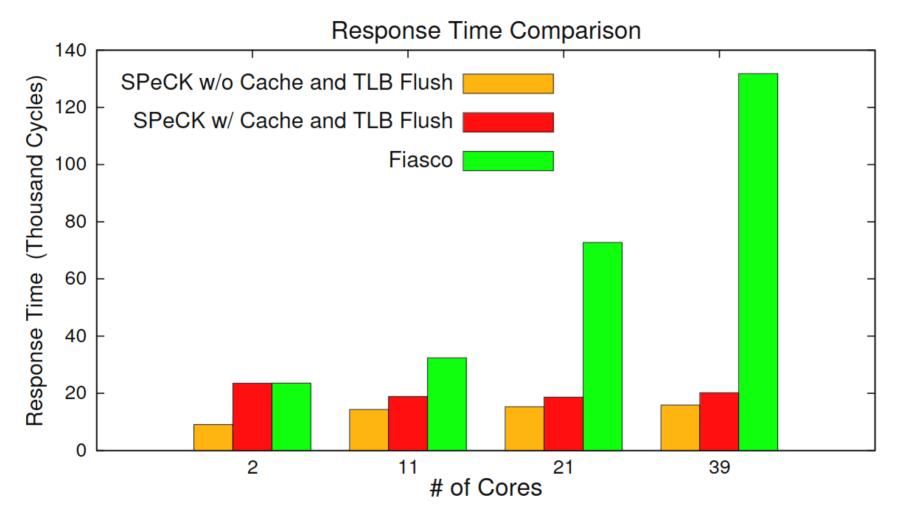
## Composite µ-kernel

- Wait-f
- No lo
  - Ato
  - Wit
  - Scal

Scalable Predictability:

Syscall worst-case on 1 core = ESyscall worst-case on N cores = cE

- Bounded execution:
  - → parallel kernel with *no locks!*

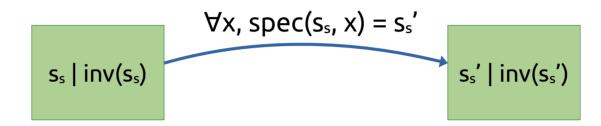


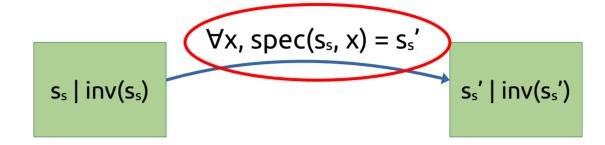
## Real-time + Scalable Predictability

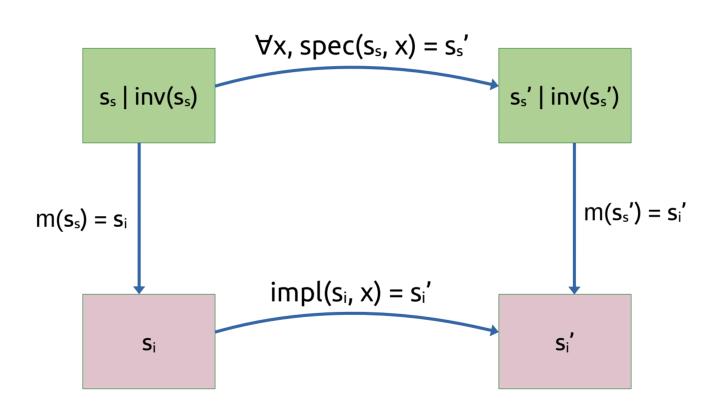
- Bounded execution
  - Constant-bounded loops
  - Bounded execution time
  - Bounded path complexity
- Wait-free
  - Syncronization: faa, non-repeating cas
  - Scalable memory reclamation
  - Bounded path complexity

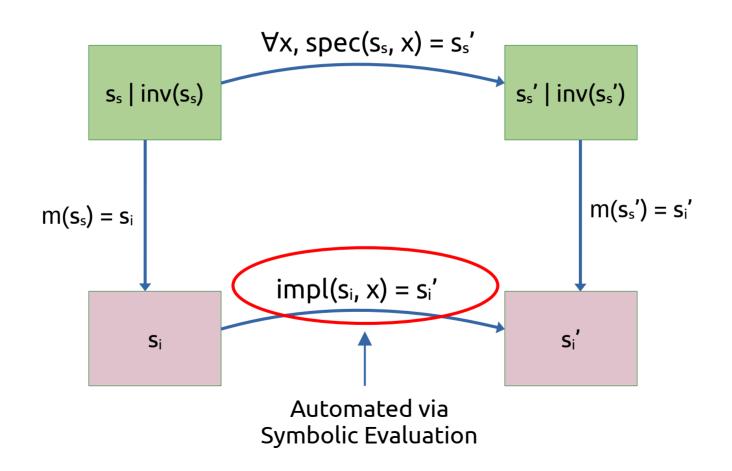
### Real-time + Scalable Predictability

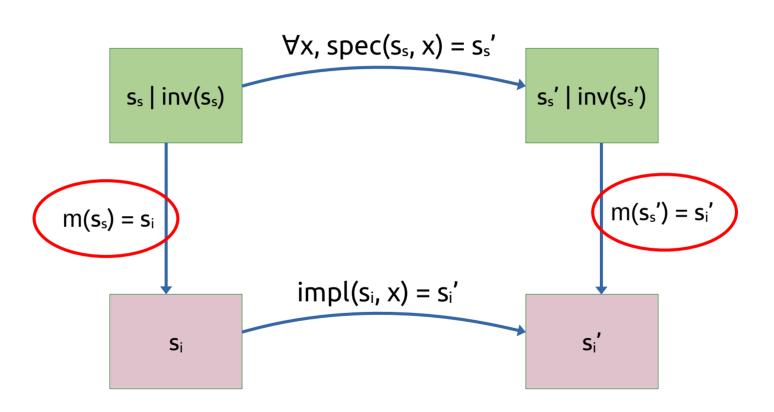
- Bounded path complexity
  - Good match for *push-button verification*
  - Example: *Serval* from U. Wash.
- Push-Button (Functional) Verification
  - Implementation (C) symbolic evaluation
  - Specification state machine transformations
  - Mapping from Implementation to Specification
  - Proof 100% via SMT solver

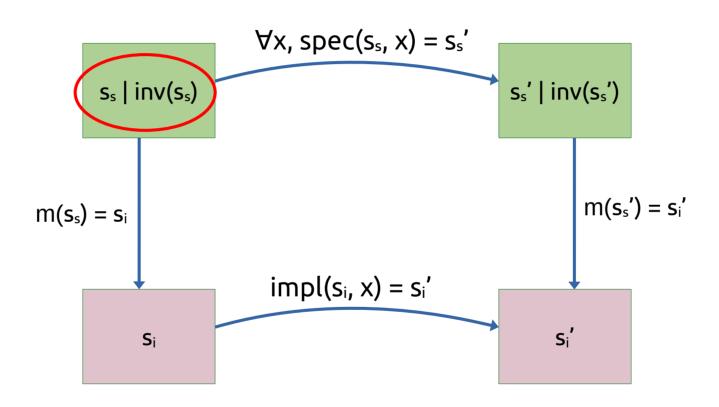


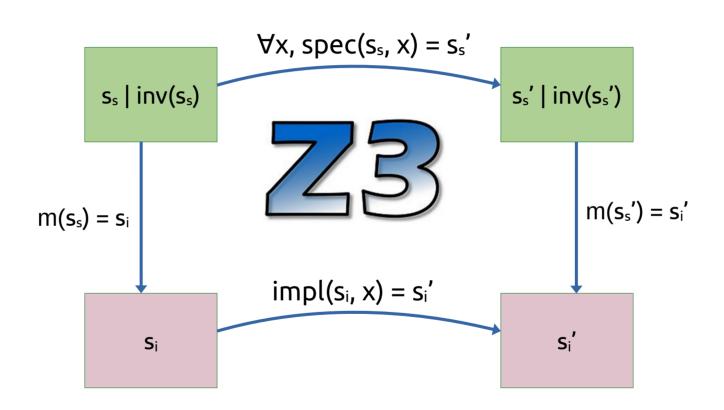












# Composite Push-Button Verification

WiP: Composite Kernel, Version 4

**Control Operations** 

Capability- & Page-tables

Typed Pages + Retyping

## Composite & Systems @ GWU

